

Algorithms

I. Asymptotic Notations & Apriori Analysis

For Micro Notes by the Student



01. State True / False

1. $100n \cdot \log n = O(n \cdot \log n)$
2. $2^{n+1} = O(2^n)$
3. $2^{2n} = O(2^n)$
4. $O < x < y$ then $n^x = O(n^y)$
5. $(n+k)^m \neq \theta(n^m)$ ($k, m > 0$)
6. $\sqrt{\log n} = O(\log \log n)$
7. $\log(n)$ is $\Omega(1/n)$
8. 2^{n^2} is $O(n!)$
9. n^2 is $O(2^{2 \log n})$
10. $a^n \neq O(n^x)$, $a > 1$, $x > 0$
11. $2^{\log_2 n^2}$ is $O(n^2)$

02. Asymptotic Comparisons

01. $f(n) = n$, $g(n) = \log n$
02. $f(n) = n^2 \log n$, $g(n) = n \cdot \log^{10} n$
03. $f(n) = n^3$, $0 < n \leq 10,000$
 $= n$, $n > 10,000$
 $g(n) = n$, $0 < n \leq 100$
 $= n^3$, $n > 100$

03. Two Packages are available for processing a Data Base having 10^x records. Package a takes a times of $10 \cdot n \cdot \log n$ while package B takes a time of $0.0001n^2$ for processing B records. Determine the smallest integer x for which Package 'A' outperforms Package 'B'.



04. Arrange the functions in increasing order of rates of growth.

01. $n^2; n \cdot \log n; n\sqrt{n}; e^n; n; 2^n; (1/n)$.

02. $2^n; n^{3/2}; n \log n; n^{\log n}$

03. $n^{(1/3)}; e^n; n^{7/4}; n \log^9 n; 1.001^n$

05. Consider the following functions from positive integers to real numbers:

$$10, \sqrt{n}, n, \log_2 n, \frac{100}{n}.$$

The CORRECT arrangement of the above functions in increasing order of asymptotic complexity is:

(a) $\log_2 n, \frac{100}{n}, 10, \sqrt{n}, n$

(b) $\frac{100}{n}, 10, \log_2 n, \sqrt{n}, n$

(c) $10, \frac{100}{n}, \sqrt{n}, \log_2 n, n$

(d) $\frac{100}{n}, \log_2 n, 10, \sqrt{n}, n$

06. Which of the following is TRUE?

f(n) is O(g(n))

g(n) is NOT O(f(n))

g(n) is O(h(n))

h(n) is O(g(n))

(a) f(n) is O(h(n))

(b) f(n) + h(n) is O(g(n)+h(n))

(c) h(n) ≠ O(f(n))

(d) f(n) · g(n) ≠ O(g(n)) · h(n)

07. An element in an Array is called Leader if it is greater than all elements to the right of it. The Time Complexity of the most efficient Algorithm to print all Leaders of the given Array of size 'n' is _____.

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Time Complexity Framework for Recursive Algorithms

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11. Algorithm What (n)

```
{
  if (n = 1) return;
  else
  {
    What (n - 1);
    B(n);
  }
}
```

12. Algorithm A(n)

```
{
  if (n = 2) return;
  else
  return  $A(\sqrt{n})$ ;
}
```

13. Algorithm Do_It (n)

```
{
  if (n = 1) return;
  else
  return (Do_It (n - 1) + Do_It (n - 1));
}
```

14. Algorithm A(n)

```
{
  if (n = 2) return;
  else
  return  $A(\sqrt{n}) + A(\sqrt{n})$ ;
}
```




```
18. void abc(char *s)
{
    if (*s != '\0')
    {
        printf("%c",*s);
        abc(s + 1);
        abc(s+1);
    }
}
```

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Running Times of Program Segments with Loops:

19. for $i \leftarrow 1$ to n
 $c = c + 1$;
20. for $i \leftarrow 1$ to n
 for $j \leftarrow 1$ to $n/2$
 $c = c + 1$;
21. for $i \leftarrow 1$ to n
 for $j \leftarrow 1$ to $n/4$
 for $k \leftarrow 1$ to n
 break;
22. for ($i = 1$; $i \leq n$; $++i$)
 for ($j = 1$; $j \leq n$; $++j$)
 for ($k = n/2$; $k \leq n$; $k += n/2$)
 $c = c + 1$;
23. $i = 1$;
 while ($i \leq n$)
 {
 $i = i * 2$;
 }
24. $i = n$;
 while ($i > 0$)
 {
 $i = i/2$;
 }

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```
25. k = 1; i = 1;
   while (k <= n)
   {
       i ++;
       k = k + i
   }
```

```
26 for (i = 1; i <= n; ++i)
    for (j = 1; j < n; j = 2 * j)
        c + c + 1;
```

```
27. m = 2n
   for (i = 1; i <= n; ++ i)
       for (j = 1; j <= m; j = 2 * j)
           c = c + 1
```

```
28. for i ← 1 to n
    for j ← i to n
```

29. $f(n) = \sum_{i=1}^n O(n)$

```
30. i = n;
   while (i > 0)
   {
       j = 1;
       while (j <= n)
       {
           j = 2 * j;
       }
       i = i/2;
   }
```

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```
31. int fun (int n)
{
  int i, j, p, q = 0 ;
  for (i = 1; i <= n ; ++ i)
  {
    p = 0;
    for (j = n; j > 1; j = j/2)
      ++ p;
    for (k = 1; k < p; k = k * 2)
      ++ q;
  }
  return (q);
}
```

```
32. for (i = 1; i <= n; ++ i)
{
  j = 1;
  while (j <= n)
    j = 2 * j;
  for (k = 1; k <= n; ++ k)
    c = c + 1
}
```

```
33. n = 22k
for (i = 1; i <= n; ++ i)
{
  j = 2;
  while (j <= n)
  {
    j = j * j;
    pf(“*”);
  }
}
```

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34. A: Array [1n] of binary;

```
f(m) =  $\theta$  (m)
count: integer;
count = 1;
for i = 1 to n
{
if (A[i] == 1) count ++;
else
{
f(count);
count = 1;
}
}
```

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35. Consider the following three functions

$$f_1 = 10^n, \quad f_2 = n^{\log n}, \quad f_3 = \sqrt{n}^n$$

Which one of the following options arranges the functions in the increasing order of asymptotic growth rate?

- (a) f_1, f_2, f_3 (b) f_3, f_2, f_1
(c) f_2, f_3, f_1 (d) f_2, f_1, f_3

36. N items are stored in a sorted doubly linked list. For a delete operation, a pointer is provided to the record to be deleted. For a decrease-key operation, a pointer is provided to the record on which the operation is to be performed. An algorithm performs the following operations on the list in this order: $\Theta(N)$ delete, $O(\log N)$ insert, $O(\log N)$ find, and $\Theta(N)$ decrease-key. What is the time complexity of all these operations put together?

- (a) $O(\log^2 N)$ (b) $O(N)$
(c) $O(N^2)$ (d) $\Theta(N^2 \log N)$



II. Divide & Conquer

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1. Assume that Merge Sort takes 30sec to Sort 64 elements in worst case. What is the approximate number of elements that can be Sorted in the Worst Case using Merge Sort using 6 minutes?

2. Consider the following recurrence relation

$$T(n) = \begin{cases} T(n/2) + T(2n/5) + 7n & \text{if } n > 0 \\ 1 & \text{if } n = 0 \end{cases}$$

Which one of the following options is correct?

- (a) $T(n) = \Theta(n \log n)$
- (b) $T(n) = \Theta((\log n)^{5/2})$
- (c) $T(n) = \Theta(n^{5/2})$
- (d) $T(n) = \Theta(n)$

3. Consider the recurrence function

$$T(n) = \begin{cases} 2T(\sqrt{n}) + 1 & n > 2 \\ 2, & 0 \leq n \leq 2 \end{cases}$$

Then $T(n)$ in terms of Θ notation is

- (a) $\Theta(\log \log n)$
- (b) $\Theta(\log n)$
- (c) $\Theta(\sqrt{n})$
- (d) $\Theta(n)$



4. For constants $a \geq 1$ and $b > 1$, consider the following recurrence defined on the non-negative integers

$$T(n) = a.T\left(\frac{n}{b}\right) + f(n)$$

Which one of the following options is correct about the recurrence $T(n)$?

- (a) If $f(n)$ is $n \log_2(n)$, then $T(n)$ is $\Theta(n \log_2(n))$
(b) If $f(n)$ is $\frac{n}{\log_2(n)}$, then $T(n)$ is $\Theta(\log_2(n))$
(c) If $f(n)$ is $O(n^{\log_b(a)-\epsilon})$ for $\epsilon > 0$, then $T(n)$ is $\Theta(n^{\log_b(a)})$
(d) If $f(n)$ is $\Theta(n^{\log_b(a)})$, then $T(n)$ is $\Theta(n^{\log_b(a)})$
5. Match the algorithms with their time complexities

Algorithm

Time complexity

- | | |
|---|-------------------------|
| (P) Towers of Hanoi with n disks | (i) $\Theta(n^2)$ |
| (Q) Binary search given n sorted numbers | (ii) $\Theta(n \log n)$ |
| (R) Heap sort given n numbers at the worst case | (iii) $\Theta(2^n)$ |
| (S) Addition of two $n \times n$ matrices | (iv) $\Theta(\log n)$ |
- (a) P (iii), Q (iv), R (i), S (ii)
(b) P (iv), Q (iii), R (i), S (ii)
(c) P (iii), Q (iv), R (ii), S (i)
(d) P (iv), Q (iii), R (ii), S (i)

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III. Greedy Method

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1. Consider the weights and values of items listed below. Note that there is only one unit of each item.

Item number	Weight (in Kgs)	Value (in Rupees)
1	10	60
2	7	28
3	4	20
4	2	24

The task is to pick a subset of these items such that their total weight is no more than 11 Kgs and their total value is maximized. Moreover, no item may be split. The total value of items picked by an optimal algorithm is denoted by V_{opt} . A greedy algorithm sorts the items by their value-to-weight ratios in descending order and packs them greedily, starting from the first item in the ordered list. The total value of items picked by the greedy algorithm is denoted by V_{greedy} . The value of $V_{opt} - V_{greedy}$ is ____.

2. We are given 9 tasks T_1, T_2, \dots, T_9 . The execution of each task requires one unit of time. We can execute one task at a time. Each task T_i has a profit P_i and a deadline d_i , Profit P_t is earned if the task is completed before the end of the Deadline.

Task	T1	T2	T3	T4	T5	T6	T7	T8	T9
Profit	15	20	30	18	18	10	23	16	25
Deadline	7	2	5	3	4	5	2	7	3

- a. Are all tasks completed in the schedule that gives maximum profit?
 - (a) All tasks are completed
 - (b) T_1 and T_6 are left out
 - (c) T_1 and T_8 are left out
 - (d) T_4 and T_6 are left out



b. What is the maximum profit earned?

- (a) 147 (b) 165 (c) 167 (d) 175

3. The characters 'a' to 'h' have the set of frequencies based on the first 8 Fibonacci numbers as follows :

a:1, b:1, c:2, d:3, e:5, f:8, g:13, h:21

A Huffman code is used to represent the characters. What is the sequence of characters corresponding to the following code?

110111100111010

- (a) fdheg (b) ecgdf
(c) dchfg (d) fehdg

4. A message is made up entirely of characters from the set $X = \{P, Q, R, S, T\}$. The table of probabilities for each of the characters is shown below:

Character	Probability
P	0.22
Q	0.34
R	0.17
S	0.19
T	0.08
Total	1.00

If a message of 100 characters over X is encoded using Huffman coding, then the expected length of the encoded message in bits is _____.

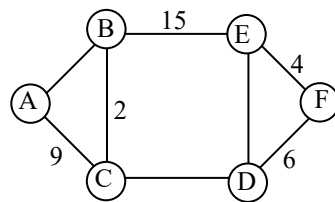
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5. Consider the string abbccddeee. Each letter in the string must be assigned a binary code satisfying the following properties:
1. For any two letters, the code assigned to one letter must not be a prefix of the code assigned to the other letter.
 2. For any two letters of the same frequency, the letter which occurs earlier in the dictionary order is assigned a code whose length is at most the length of the code assigned to the other letter.
- Among the set of all binary code assignments which satisfy the above two properties, what is the minimum length of the encoded string?
- (a) 25 (b) 23 (c) 21 (d) 30

6. Consider a Graph whose vertices are points in a plane with integer co-ordinates (x, y) where $1 \leq x \leq n, 1 \leq y \leq n, n > 2$ is an integer. 2 vertices $\langle x_1, y_1 \rangle$ & $\langle x_2, y_2 \rangle$ are adjacent iff $|x_1 - x_2| \leq 1$ & $|y_1 - y_2| \leq 1$. The cost of such an edge is given by the distance between them. Compute the weight of min cost Spanning Tree of such graph for a value of n .
7. Consider the following Graph whose Minimum Cost Spanning Tree marked with edge values has a weight of 36. Minimum possible sum of all edges of the graph G is _____. (Assume that all edges have distinct cost).



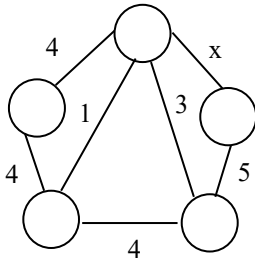
8. Consider a graph with 'n' vertices $n > 2$. The vertices are numbered V_1 to V_n . Two vertices V_i & V_j are adjacent iff $0 < |i - j| \leq 2$. The weight of such an edge is $i + j$. The weight of minimum cost Spanning Tree of such a graph for a value of n is _____.
9. Consider a complete weighted Graph with n vertices numbered V_1 to V_n . Two vertices V_i & V_j having edge between them has a cost value of $2|i - j|$. The weight of minimum cost Spanning Tree of such a graph is _____.

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10. Let G be a complete undirected graph with 4 vertices and edge weights are $\{1, 2, 3, 4, 5, 6\}$. The maximum possible weight that a minimum weight Spanning Tree can have is _____.
11. Let G be a connected undirected graph of 100 vertices and 300 edges. The weight of a minimum spanning tree of G is 500. When the weight of each edge of G is increased by five, the weight of a minimum spanning tree becomes _____.
12. Consider the following undirected graph G :



- Choose a value for x that will maximize the number of minimum weight spanning trees (MWSTs) of G . The number of MWSTs of G for this value of x is _____.
13. Let w be the minimum weight among all edge weights in an undirected connected graph. Let 'e' be a specific edge of weight 'w'. Which of the following is False?
- There is a minimum Spanning Tree containing 'e' always.
 - Every minimum Spanning Tree has an edge of weight 'w'.
 - 'e' is present in every minimum Spanning Tree.
 - If 'e' is not present in a minimum Spanning Tree named 'T' then there will be a cycle formed by adding 'e' to T.
14. $G = (V, E)$ is an undirected simple graph in which each edge has a distinct weight, and e is a particular edge of G . Which of the following statements about the minimum spanning trees (MSTs) of G is/are **TRUE**?

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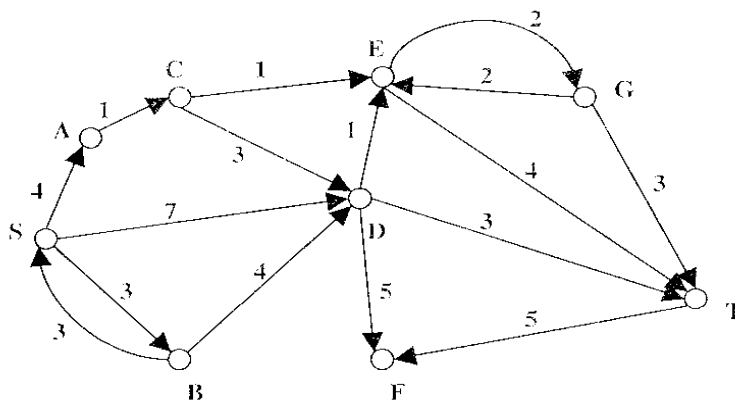


I. If e is the lightest edge of some cycle in G , then every MST of G includes e

II. If e is the heaviest edge of some cycle in G , then every MST of G excludes e

- (a) I only (b) II only
(c) both I and II (d) neither I nor II

15. Applying Dijkstra's Algorithm over the given Graph, Which path is reported from 'S' to 'T';



16. Let G be a weighted connected undirected graph with distinct positive edge weights. If every edge weight is increased by the same value, then which of the following statements is/are true?

1. Minimum spanning Tree of the graph does not change.
2. Shortest path between any pair of vertices does not change.

17. Consider the weighted undirected graph with 4 vertices, where the weight of edge $\{i,j\}$ is given by the entry W_{ij} in the matrix W .

$$W = \begin{bmatrix} 0 & 2 & 8 & 5 \\ 2 & 0 & 5 & 8 \\ 8 & 5 & 0 & x \\ 5 & 8 & x & 0 \end{bmatrix}$$

The largest possible integer value of x , for which at least one shortest path between some pair of vertices will contain the edge with weight x is _____.

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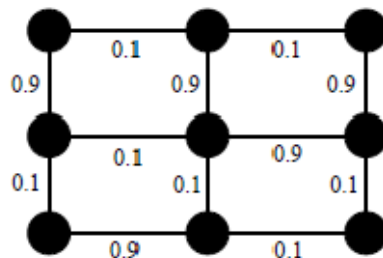
18. Let $G = (V, E)$ be any connected undirected edge-weighted graph. The weights of the edges in E are positive and distinct. Consider the following statements:

- (I) Minimum Spanning Tree of G is always unique.
- (II) Shortest path between any two vertices of G is always unique.

Which of the above statements is/are necessarily true?

- (a) (I) only
- (b) (II) only
- (c) both (I) and (II)
- (d) neither (I) nor (II)

19. Consider the following undirected graph with edge weights as shown



The number of minimum-weight spanning trees of the graph is _____

20. Let $G = (V, E)$ be any connected undirected edge-weighted graph. The weights of the edges in E are positive and distinct. Consider the following statements:

- (I) Minimum Spanning Tree of G is always unique.
- (II) Shortest path between any two vertices of G is always unique.

Which of the above statements is/are necessarily true?

- (a) (I) only
- (b) (II) only
- (c) both (I) and (II)
- (d) neither (I) nor (II)

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21. Let G be a connected undirected weighted graph. Consider the following two statements.

S₁: There exists a minimum weight edge in G which is present in every minimum spanning tree of G

S₂: If every edge in G has distinct weight, then G has a unique minimum spanning tree.

Which one of the following options is correct?

- (a) Both S₁ and S₂ are true
- (b) S₁ is true and S₂ is false
- (c) S₁ is false and S₂ is true
- (d) Both S₁ and S₂ are false.

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IV. Dynamic Programming

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1. Consider the following table:

Algorithms	Design Paradigms
(P) Kruskal	(i) Divide and Conquer
(Q) Quick sort	(ii) Greedy
(R) Floyd-Warshall	(iii) Dynamic Programming

Match the algorithms to the design paradigms they are based on.

- (a) (P) ↔ (ii), (Q) ↔ (iii), (R) ↔ (i)
(b) (P) ↔ (iii), (Q) ↔ (i), (R) ↔ (ii)
(c) (P) ↔ (ii), (Q) ↔ (i), (R) ↔ (iii)
(d) (P) ↔ (i), (Q) ↔ (ii), (R) ↔ (iii)
2. Which one of the following algorithm design techniques is used in solving Sum of Subsets problem?
- (a) Dynamic programming
(b) Backtracking
(c) Greedy Technique
(d) Divide and Conquer

< Linked Question >

The subset-sum problem is defined as follows. Given a set of n positive integers, $S = \{a_1, a_2, a_3, \dots, a_n\}$, and positive integer W , is there a subset of S whose elements sum to W ? A dynamic program for solving this problem uses a 2-dimensional Boolean array, X , with n rows and $W+1$ columns $X[i, j], 1 \leq i \leq n, 0 \leq j \leq W$, is TRUE if and only if there is a subset of $\{a_1, a_2, a_3, \dots, a_n\}$ whose elements sum to j .



3. Which of the following is valid for $2 \leq i \leq n$ and $a_i \leq j \leq W$?
- (a) $X[i, j] = X[i-1, j] \sqcup X[i, j-a_i]$
 - (b) $X[i, j] = X[i-1, j] \sqcup X[i-1, j-a_i]$
 - (c) $X[i, j] = X[i-1, j] \sqcup X[i, j-a_i]$
 - (d) $X[i, j] = X[i-1, j] \sqcup X[i-1, j-a_i]$
4. Which entry of the array X, if **TRUE**, implies that there is a subset whose elements sum to W?
- (a) $X[1, w]$ (b) $X[n, 0]$
 - (c) $X[n, W]$ (d) $X[n-1, n]$
5. Consider two strings $A = "qpqrr"$ and $B = "pqprrrp"$. Let x be the length of the longest common subsequence (not necessarily contiguous) between A and B and let y be the number of such longest common subsequences between A and B. Then $x + 10y = \underline{\hspace{2cm}}$.

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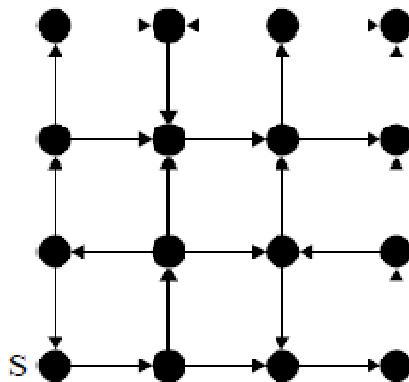


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V. Graph Techniques & Components

1. A DFS is performed on DAG. Which of the following is true for all edges (u, v) in the graph?
 - (a) $d[u] < d[v]$
 - (b) $d[u] < f[v]$
 - (c) $f[u] < f[v]$
 - (d) $f[u] > f[v]$
2. Consider a DFT of an undirected graph having 'n' vertices. In the traversal, k edges are marked as Tree edges then the number of connected components in the graph is given by
 - (a) k
 - (b) k + 1
 - (c) n - k
 - (d) n - k - 1
3. Consider the following directed graph

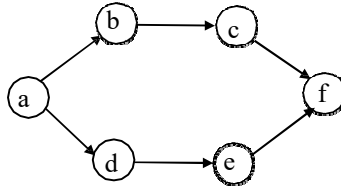


Which of the following is/are correct about the graph?

- (a) The graph does not have a strongly connected component.
- (b) A depth-first traversal starting at vertex S classifies three directed edges as back edges.
- (c) For each pair of vertices u and v, there is a directed path from u to v.
- (d) The graph does not have a topological order.



4. Consider the following directed graph:



The number of different topological orderings of the vertices of the graph is _____ .

5. An articulation point in a connected graph is a vertex such that removing the vertex and its incident edges disconnects the graph into two or more connected components.

Let T be a DFS tree obtained by doing DFS in a connected undirected graph G. Which of the following options is/are correct?

- (a) Root of T can never be an articulation point in G
- (b) If u is an articulation point in G such that x is an ancestor of u in T and y is a descendant of u in T, then all paths from x to y in G must pass through u.
- (c) A leaf of T can be an articulation point in G
- (d) Root of T is an articulation point in G if and only if it has 2 or more children.

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VI. Heaps

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- Valid binary Max-Heap
(a) $\langle 25, 12, 16, 13, 10, 8, 14 \rangle$ (b) $\langle 25, 14, 16, 13, 10, 8, 12 \rangle$
(c) $\langle 25, 14, 13, 16, 10, 8, 12 \rangle$ (d) $\langle 25, 14, 12, 13, 10, 8, 16 \rangle$
- Valid 3-ary maximum Heap Array representation
(a) $\langle 1, 3, 5, 6, 8, 9 \rangle$ (b) $\langle 9, 6, 3, 1, 8, 5 \rangle$
(c) $\langle 9, 3, 6, 8, 5, 1 \rangle$ (d) $\langle 9, 5, 6, 8, 3, 1 \rangle$
- To the valid Heap of Q55 insert elements $\langle 7 \ 2 \ 10 \ 4 \rangle$. Indicate the resultant Heap in Array.
- Level order traversal of a binary max Heap generates: $\langle 10, 8, 5, 3, 2 \rangle$
Insert: $\langle 1 \ \& \ 7 \rangle$; What is the resultant Level order Traversal?
- In a binary max-Heap with n elements, the smallest element can be found in time of _____.
- Given binary Heap with 'n' elements & it is required to insert 'n' more elements not necessarily one after another into this Heap. Total time required for this operation is:
(a) $O(n^2)$ (b) $n \log n$ (c) n (d) $n^2 \log n$
- Given binary Heap in Array with the smallest at the root, the 7th smallest element can be found in time complexity of _____.
- Consider binary Heap in an Array with n elements it is desired to insert an element into the Heap if a binary search is performed along the path from newly inserted element to the root then the no. of comparison made is order of _____.
- The approximate no. of elements that can be Sorted in $O(\log n)$ time using Heap Sort is _____.



15. Consider a complete binary tree where the left and right subtrees of the root are max-heaps. The lower bound for the number of operations to convert the tree to a heap is

- (a) $\Omega(\log n)$
- (b) $\Omega(n)$
- (c) $\Omega(n \log n)$
- (d) $\Omega(n^2)$

16. Consider a max heap, represented by the array:

40, 30, 20, 10, 15, 16, 17, 8, 4.

Array index	1	2	3	4	5	6	7	8	9
Value	40	30	20	10	15	16	17	8	4

Now consider that a value 35 is inserted into this heap. After insertion, the new heap is

- (a) 40, 30, 20, 10, 15, 16, 17, 8, 4, 35
- (b) 40, 35, 20, 10, 30, 16, 17, 8, 4, 15
- (c) 40, 30, 20, 10, 35, 16, 17, 8, 4, 15
- (d) 40, 35, 20, 10, 15, 16, 17, 8, 4, 30

17. The number of possible min-heaps containing each value from {1, 2, 3, 4, 5, 6, 7} exactly once is _____

18. Consider the following statements:

- I. The smallest element in a max-heap is always at a leaf node.
- II. The second largest element in a max-heap is always a child of the root node.
- III. A max-heap can be constructed from a binary search tree in (n) time.
- IV. A binary search tree can be constructed from a max-heap in (n) time.

Which of the above statements are TRUE?

- (a) I, III and IV
- (b) II, III and IV
- (c) I, II and III
- (d) I, II and IV

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VII. Sorting Methods

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- Which of the following Sorting algorithms has lowest worst-case complexity?
 - Bubble Sort
 - Merge Sort
 - Quick Sort
 - Selection Sort
- Which of the following in place Sorting algorithm needs minimum number of swaps?
 - Selection Sort
 - Insertion Sort
 - Heap Sort
 - Quick Sort
- What would be the worst case complexity of Insertion Sort if the inputs are restricted to permutation of 1 to n with at most 'n' Inversions?
- Let 'S' be a Sorted Array of 'n' integers and T(n) denote the time taken for the most efficient algorithm to determine if there are 2 elements in the Array with the sum <1000.
 - T(n) is O(1)
 - $n \leq T(n) \leq n \log n$
 - $T(n) = {}^n C_2$
 - $n \log n \leq T(n) = {}^n C_2$
- The traditional Insertion Sort to Sort an Array can be of n elements uses linear search to identify the position where an element is to be inserted into already Sorted part of the Array, if instead binary search is used to identify the position of newly inserted element then the worst case complexity will be order of _____.
- In using Quick Sort suppose the central element of the Array is always chosen as the Pivot then the worst case complexity of the Quick Sort may be _____.
- The Median on Array of size n can be found in O(n) time. If Median is selected as Pivot, then the worst case complexity of Quick Sort is _____.
- In applying Quick Sort to an unsorted list if (n/4) the element is selected as Pivot then the Time Complexity of Quick Sort will be _____.



9. Consider the Quicksort algorithm. Suppose there is a procedure for finding a pivot element which splits the list into two sublists each of which contains at least one-fifth of the elements. Let $T(n)$ be the number of comparisons required to sort n elements. Then
- (a) $T(n) \leq 2T(n/5)+n$ (b) $T(n) \leq T(n/5)+T(4n/5)+n$
(c) $T(n) \leq 2T(4n/5)+n$ (d) $T(n) \leq 2T(n/2)+n$
10. Which one of the following in place sorting algorithms needs the minimum number of swaps?
- (a) Quick sort (b) Insertion sort
(c) Selection sort (d) Heap sort
11. The worst case running times of Insertion sort, Merge sort and Quick sort, respectively, are:
- (a) $\Theta(n \log n)$, $\Theta(n \log n)$, and $\Theta(n^2)$
(b) $\Theta(n^2)$, $\Theta(n^2)$, and $\Theta(n \log n)$
(c) $\Theta(n^2)$, $\Theta(n \log n)$, and $\Theta(n \log n)$
(d) $\Theta(n^2)$, $\Theta(n \log n)$, and $\Theta(n^2)$
12. If one uses straight two-way merge sort algorithm to sort the following elements in ascending order:
20, 47, 15, 8, 9, 4, 40, 30, 12, 17
then the order of these elements after second pass of the algorithm is:
- (a) 8, 9, 15, 20, 47, 4, 12, 17, 30, 40
(b) 8, 15, 20, 47, 4, 9, 30, 40, 12, 17
(c) 15, 20, 47, 4, 8, 9, 12, 30, 40, 17
(d) 4, 8, 9, 15, 20, 47, 12, 17, 30, 40
13. You have n lists, each consisting of m integers sorted in ascending order. Merging these lists into a single sorted list will take time:
- (a) $O(nm \log m)$ (b) $O(mn \log m)$
(c) $O(m + n)$ (d) $O(mn)$

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14. If we use Radix Sort to sort n integers in the range $(n^{k/2}, n^k]$, for some $k > 0$ which is independent of n , the time taken would be
- (a) $\Theta(n)$ (b) $\Theta(kn)$
(c) $\Theta(n \log n)$ (d) $\Theta(n^2)$
15. The worst case running times of Insertion sort, Merge sort and Quick sort, respectively are:
- (a) $\Theta(n \log n)$, $\Theta(n \log n)$ and $\Theta(n^2)$
(b) $\Theta(n^2)$, $\Theta(n^2)$ and $\Theta(n \log n)$
(c) $\Theta(n^2)$, $\Theta(n \log n)$ and $\Theta(n \log n)$
(d) $\Theta(n^2)$, $\Theta(n \log n)$ and $\Theta(n^2)$
16. Let P be a quick sort program to sort numbers in ascending order. Let t_1 and t_2 be the time taken by the program for the inputs $[1\ 2\ 3\ 4]$ and $[5\ 4\ 3\ 2\ 1]$, respectively. Which of the following holds?
- (a) $t_1 = t_2$ (b) $t_1 > t_2$ (c) $t_1 < t_2$ (d) $t_1 = t_2 = 5 \log 5$
17. Let P be a Quick Sort Program to sort numbers in ascending order using the first element as pivot. Let t_1 and t_2 be the number of comparisons made by P for the inputs $\{1, 2, 3, 4, 5\}$ and $\{4, 1, 5, 3, 2\}$ respectively. Which one of the following holds?
- (a) $t_1 = 5$ (b) $t_1 < t_2$ (c) $t_1 > t_2$ (d) $t_1 = t_2$
18. Quick-sort is run on two inputs shown below to sort in ascending order taking first element as pivot
- i. $1, 2, 3, \dots, n$
ii. $n, n-1, n-2, \dots, 2, 1$
- Let C_1 and C_2 be the number of comparisons made for the inputs (i) and (ii) respectively. Then,
- (a) $C_1 < C_2$ (b) $C_1 > C_2$
(c) $C_1 = C_2$ (d) We cannot say anything for arbitrary n

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19. Following algorithm(s) can be used to sort n in the range $[1..n^3]$ in $O(n)$ time

- (a) Heap sort (b) Quick sort
(c) Merge sort (d) Radix sort

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20. For merging two sorted lists of sizes m and n into a sorted list of size $m+n$, we require comparisons of

- (a) $O(m)$ (b) $O(n)$ (c) $O(m+n)$ (d) $O(\log m + \log n)$

21. Give the correct matching for the following pairs:

- (A) $O(\log n)$ (P) Selection sort
(B) $O(n)$ (Q) Insertion sort
(C) $O(n \log n)$ (R) Binary search
(D) $O(n^2)$ (S) Merge sort

22. Assume that the algorithms considered here sort the input sequences in ascending order. If the input is already in ascending order, which of the following are **TRUE** ?

- I. Quicksort runs in $\Theta(n^2)$ time
II. Bubblesort runs in $\Theta(n^2)$ time
III. Mergesort runs in $\Theta(n)$ time
IV. Insertion sort runs in $\Theta(n)$ time
(a) I and II only (b) I and III only
(c) II and IV only (d) I and IV only

23. An array of 25 distinct elements is to be sorted using quicksort. Assume that the pivot element is chosen uniformly at random. The probability that the pivot element gets placed in the worst possible location in the first round of partitioning (rounded off to 2 decimal places) is _____.